

Concise representation of hypergraph minimal transversals : approach and application on the dependency inference problem

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Motivations

- Hypergraph minimal transversals (MTs) are useful in many computer science applications [Hagen, 2008]:
 - ▶ Data mining, AI, Databases, Semantic web, etc.
- Several algorithms have been proposed in the literature.
 - ▶ Berge¹, Fredman et al.², Kavvadias et al.³, MTMINER⁴, etc.
 - ▶ Number of MTs may be **exponential** relative to hypergraph size.
- Our approach : An **information lossless concise** representation of MTs.
 - ▶ Notion of irredundant hypergraph.
 - ▶ Proposition of a new algorithm, IRRED-ENGINE, for the extraction of all the MTs.

¹[Berge, 1989]

²[Fredman, 1996]

³[Kavvadias et Stavropoulos, 2005]

⁴[Hébert, 2007]

Outline

- 1 Key notions
- 2 Irredundant MTs
- 3 Approach and algorithm
- 4 Experimentations
- 5 Application : dependency inference problem
- 6 Conclusion

Outline

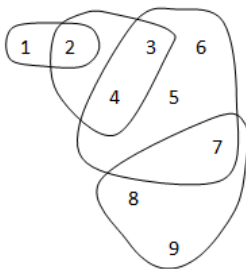
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Hypergraph and minimal transversal

- Let $H = (\mathcal{X}, \xi)$ such as :
 - ▶ $\mathcal{X} = \{x_1, x_2, \dots, x_n\}$: a finite set of vertices.
 - ▶ $\xi = \{e_1, e_2, \dots, e_m\}$: a set of subsets of \mathcal{X} .
 - $H = (\mathcal{X}, \xi)$ is a hypergraph if [Berge, 1989] :
 - ▶ $e_i \neq \emptyset, i \in \{1, \dots, m\}$
 - ▶ $\bigcup_{i=1, \dots, m} e_i = \mathcal{X}$
 - $T \subset \mathcal{X}$ is a transversal of H if $T \cap e_i \neq \emptyset \forall i = 1, \dots, m$.
 - T is called minimal if $\nexists T_1 \subset T$ s.t. T_1 is a transversal.
- The set of MTs is denoted \mathcal{M}_H .

■ Let $H = (\mathcal{X}, \xi)$ such as :

- ▶ $\mathcal{X} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- ▶ $\xi = \{e_1, e_2, e_3, e_4\}$



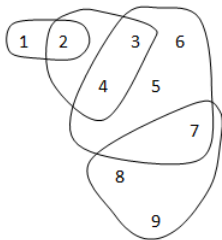
$$e_1 = \{1, 2\}$$

$$e_2 = \{2, 3, 4\}$$

$$e_3 = \{3, 4, 5, 6, 7\}$$

$$e_4 = \{7, 8, 9\}$$

■ A hypergraph :



■ 15 MTs :

- ▶ $\{2, 7\} - \{1, 3, 7\}$
- ▶ $\{1, 4, 7\} - \{1, 3, 8\}$
- ▶ $\{1, 3, 9\} - \{1, 4, 8\}$
- ▶ $\{1, 4, 9\} - \{2, 3, 8\}$
- ▶ $\{2, 4, 8\} - \{2, 4, 9\}$
- ▶ $\{2, 3, 9\} - \{2, 5, 8\}$
- ▶ $\{2, 6, 8\} - \{2, 5, 9\}$
- ▶ $\{2, 6, 9\}$

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Definitions and notations

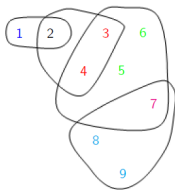
- **EXTENT** (x): extension of a vertex $x \in \mathcal{X}$.
 - ▶ The set of hyperedges to which belongs x .
- **GENERALIZED NODE** : set of vertices which have the same extension.
- **REPRESENTATIVE** : the first element of a generalized node, whenever its elements are sorted w.r.t lexicographic order.
 - ▶ Each **GENERALIZED NODE** has only one **REPRESENTATIVE**.

Notion of irredundance

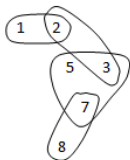
- Irredundant hypergraph : $H' = (\mathcal{X}', \xi')$ of $H = (\mathcal{X}, \xi)$.
 - ▶ $\mathcal{X}' \subseteq \mathcal{X}$: the set of the representatives of the generalized nodes associated to the vertices of H .
 - ▶ ξ' : the set of hyperedges of H deprived of the elements of $\mathcal{X} - \mathcal{X}'$
- Irredundant minimal transversals ($\mathcal{M}_{H'}$)
 - ▶ $TM_{H'} \Leftrightarrow TMirr_H$

Study case

Initial hypergraph H



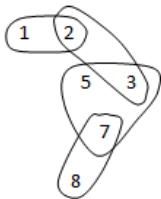
Irredundant hypergraph H'



Vertex	Extension	G.N (Rep)
1	$\{e_1\}$	$S_1(1)$
2	$\{e_1, e_2\}$	$S_2(2)$
3	$\{e_2, e_3\}$	$S_3(3)$
4	$\{e_2, e_3\}$	$S_3(3)$
5	$\{e_3\}$	$S_4(5)$
6	$\{e_3\}$	$S_4(5)$
7	$\{e_3, e_4\}$	$S_5(7)$
8	$\{e_4\}$	$S_6(8)$
9	$\{e_4\}$	$S_6(8)$

Study case

■ Hypergraph H'



■ $\mathcal{M}_{H'}$:

- ▶ {2, 7}
- ▶ {1, 3, 7}
- ▶ {1, 3, 8}
- ▶ {2, 3, 8}
- ▶ {2, 5, 8}

■ $|\mathcal{M}_{H'}| = 5$

■ $|\mathcal{M}_H| = 15$

Key notions

Irredundant MTs

Approach and algorithm

Experimentations

Application : dependency inference problem

Conclusion

The irred-engine algorithm

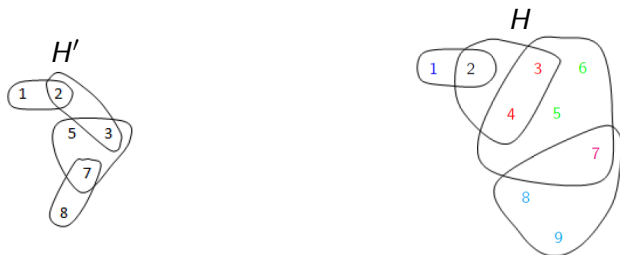
Study case

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Steps

- Input : an hypergraph H
 - Output : The set of MTs of H
- ① Computation of all generalized nodes of H .
 - ② Generation of H' .
 - ③ Extraction of MTs of H' ($\mathcal{M}_{H'}$) using the representative set.
 - ④ Derive faithfully all the MTs of H from $\mathcal{M}_{H'}$ using a simple substitution process.



$\mathcal{M}_{H'}$	\mathcal{M}_H
$\{2, 7\}$	$\{2, 7\}$
$\{1, 3, 7\}$	$\{1, 3, 7\}, \{1, 4, 7\}$
$\{1, 3, 8\}$	$\{1, 3, 8\}, \{1, 4, 8\}, \{1, 4, 9\}, \{1, 3, 9\}$
$\{2, 3, 8\}$	$\{2, 3, 8\}, \{2, 4, 8\}, \{2, 4, 9\}, \{2, 3, 9\}$
$\{2, 5, 8\}$	$\{2, 5, 8\}, \{2, 5, 9\}, \{2, 6, 8\}, \{2, 6, 9\}$

MTs of H' and MTs of H

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Datasets

Experiments were carried out on several types of hypergraphs :

- The first one are generated from the datasets " *Accidents*" and " *Connect-4*"
- The second one are a randomly hypergraphs generated using the " *random hypergraph generator*" (Boros et al.,2003).

Algorithms considered

- IRRED-ENGINE
- MMCS (Murakami et al., 2013)
- KS (Kavvadias et al., 2005)

Experimentations

- Compactness rate $\theta = 1 - |\mathcal{M}_{H'}| / |\mathcal{M}_H|$.
 - ▶ The percentage of MTs that can be built from the irredundant MTs, without loss of information.

	$ \mathcal{M}_H $	$ \mathcal{M}_{H'} $	θ
Accidents1	1.961	1.866	4.84%
Accidents2	17.486	17.199	1.64%
Connect-Win	4.587.967	4.423.837	3.57%

Accidents and Connect Hypergraphs statistics

	Accidents1	Accidents2	Connect-Win
MMCS(H)	0.301	2.787	88.491
KS(H)	8.620	-	-
MMCS(H')	0.001	2.366	84.015
+			
GET-ALL-MT($\mathcal{M}_{H'}$)	0.035	0.010	7.291
=			
IRRED-ENGINE (H)	0.036	2.376	91.306

Computation time on Accidents and Connect Hypergraphs (in seconds)

	$ \mathcal{X} $	$ \xi $	$ Gen_nodes $	$\tau(H)$
$H1$	95	51	65	8
$H2$	99	101	64	9
$H3$	117	20005	68	4
$H4$	60	20	20	20

	$ \mathcal{M}_H $	$ \mathcal{M}_{H'} $	θ
$H1$	832.564.740	358.392	99.95%
$H2$	265.765.380	189.444	99.92%
$H3$	1693	1250	26.17%
$H4$	3^{20}	1	99.99%

Matching and randomly generated hypergraphs characteristics

	$H1$	$H2$	$H3$	$H4$
MMCS(H)	2083.792	766.19	124.692	245.610
KS(H)	-	-	1041.29	18.550
MMCS(H')	2.740	0.37	87.54	0.001
+				
GET-ALL-MT($\mathcal{M}_{H'}$)	11.696	1.53	0.031	19.806
=				
IRRED-ENGINE (H)	14.436	1.90	87.571	19.807

Computation time on matching and randomly generated hypergraphs (in seconds)

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Main challenge

- Illustrate the relevance of our concise representation of hypergraph MTs
- Improve the dependency inference problem :
 - ▶ Computing a compact cover of FDs that faithfully represents the set of all FD holding in a relation r .

Functional dependency

Given a relation schema R , a FD over R is an expression $X \rightarrow Y$, where $X, Y \subseteq R$. A FD holds in r , denoted $r \models X \rightarrow Y$, where r is a relation over R , if all tuples $u, v \in R$ with $u[X] = v[X]$ satisfy also $u[Y] = v[Y]$.

Related work

- Several approach proposed in the literature classified into two categories : top-down algorithms and bottom-up algorithms.
- Bottom-up algorithms compute a subset of FD , called the cover of the fd and denoted $COVER(DEP(r))$.
- Notions of agree sets, maximal sets, disagree sets, necessary sets, etc.
- The minimal cover problem can be solved by computing the hypergraph MTs.
- **Approach** : Computing an irredundant minimal cover $COVER'(D_r)$
 - ▶ To illustrate our proposition, we consider $DEP-MINER$ [Lopes et al., 2000]

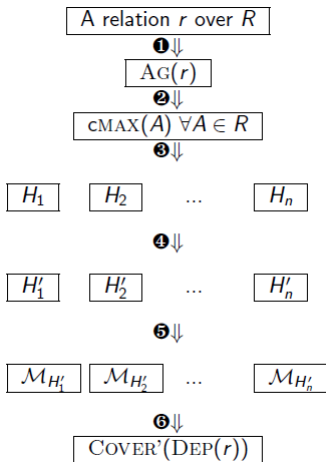


Figure: Methodology description

Study case

	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>
t_1	1	100	1	2	50
t_2	4	101	1	2	50
t_3	1	102	2	2	70
t_4	1	200	1	2	50
t_5	2	101	3	3	100
t_6	2	200	1	3	70
t_7	1	100	3	2	50

A relation r

$BE \rightarrow A$	$A \rightarrow D$	$AB \rightarrow E$
$BD \rightarrow A$	$CE \rightarrow D$	$BD \rightarrow E$
	$BE \rightarrow D$	$AC \rightarrow E$
		$CD \rightarrow E$

COVER(D_r)

Study case

	A	B	C	D	E
t_1	1	100	1	2	50
t_2	4	101	1	2	50
t_3	1	102	2	2	70
t_4	1	200	1	2	50
t_5	2	101	3	3	100
t_6	2	200	1	3	70
t_7	1	100	3	2	50

A relation r

$BE \rightarrow A$	$A \rightarrow D$	$AB \rightarrow E$
$BD \rightarrow A$	$CE \rightarrow D$	$BD \rightarrow E$
	$BE \rightarrow D$	$AC \rightarrow E$
		$CD \rightarrow E$

COVER(D_r)

$BD \rightarrow A$
$A \rightarrow D$
$BE \rightarrow D$
$AB \rightarrow E$

COVER'(D_r)

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Conclusion

- We propose an approach to improve hypergraph minimal transversals computation.
 - ▶ Concise representation of the set of MTs.
 - ▶ Notions of irredundant hypergraph and irredundant MT.
- We introduce a new algorithm, called IRRED-ENGINE which extracts all the MTs from a subset.
- Experiments, carried out on several types of hypergraphs, showed very interesting results evaluated through a compactness measure.
- Application on the dependency inference problem.

Thanks for your attention...